Microsoft Research Faculty Summit 2012
ADVANCING THE STATE OF THE ART
Operating System Architecture For Parallel And Distributed Computing

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Most of it dates back to the 60’s and 70’s

It locks us in to old ways of doing things
Using a few identical processor cores
Using the OS kernel to assign cores to threads (“kernel threads”)
Using OS kernel operations to synchronize threads (typically by locking)
Using interrupts to synchronize threads with devices
Using traps to let threads invoke operating system services
Using a variety of heuristics to allocate and schedule resources

These approaches don’t cut it any more...
Hardware Has Changed

**Available battery life limits user experiences**
There isn’t enough energy to run all of the hardware all of the time.
How can the state of battery charge be taken into account?

**Processor cores are increasingly heterogeneous**
Application software developers are trying to exploit them.
Can every core type access the OS services it needs?

**New sensors and interface devices are emerging**
Some of these require substantial or timely resource allocation.
How are these resource allocations decided and implemented by the OS?
Parallelism varies in both form and quantity
Even different phases of a single application can exhibit such diversity. How can the OS adapt to varying application parallelism?

Media and games require responsiveness
Different platforms need different resource allocations even for the same function. How can appropriate allocation be done, even given multiple simultaneous instances?

Applications are built from distributed services
These services provide search, data access, connectivity, computation, and the like. How well does application performance compose from the services it comprises?
Vulnerability Has Changed

Attacks frequently target drivers and system services.
These are increasingly large and complex software subsystems.
Must the OS share most of its resources with its device drivers and system services?

Distributed applications often connect via the internet
Messages among application components need authentication and privacy.
How do we ensure trust for both ends of these internet connections?

Internet security is an ongoing multi-player contest
The forces of good must keep on improving to keep on winning.
How can we defend against complete OS takeovers by the opposition?
Unconventional Approaches

Remove drivers from the kernel, kernel threads, manage resources with capabilities

Exokernels: Xok (1997)
Unprivileged application use of hardware, kernel threads, capability-based memory sharing

Message oriented; S: safe pointers and kernel threads, B: capabilities and user threads.

My own work: TeraOS (1992)
Reduced driver privilege, user threads, no interrupts
What Tomorrow’s Client OS Might Look Like
Continuously minimize the total penalty of the system

(subject to the total resources available)

Power And Battery Energy Management

Treat the battery as a competitor for resources

The battery is associated with a distinct process, Process\(0\).
All slack resources are assigned to Process\(0\) which tries to power them down.
As resources power down the system power diminishes in a convex manner.
System power thus acts as the “Runtime” for Process\(0\).
Penalty\(_0\) lets the system or user describe the relative importance of remaining battery life.
Other Stuff Currently Underway

Dynamic Root of Trust Measurement (DRTM)
Measured from a small and secure Trusted Platform Module (TPM)

Asynchronous interfaces in Microsoft software
Asynchronous patterns in C#/F#/VB; reactive framework; more coming soon

C++ AMP for programming GPUs or CPUs
More work to do to make the two work together smoothly
Conclusions

Operating Systems are changing (at long last).
Opinions above are mine but maybe not Microsoft’s.
Your questions and comments are welcome.