AN EMBEDDED DCT APPROACH TO PROGRESSIVE IMAGE COMPRESSION

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ABSTRACT
Motivated by Shapiro's embedded zerotree wavelet coding (EZW) and Taubman and Zakhor’s layered zero coding (LZC), we propose a layered DCT image compression scheme, which generates an embedded bit stream for DCT coefficients according to their importance. The new method allows progressive image transmission and simplifies the rate-control problem. In addition to these functionalities, it provides a substantial rate-distortion improvement over the JPEG standard when the bit rates become low. For example, we observe a bit rate reduction with respect to the JPEG Huffman and arithmetic coders by about 60% and 20%, respectively, for a bit rate around 0.1bpp.

Keywords: Layered coding, progressive transmission, arithmetic coder, JPEG

1. INTRODUCTION
The JPEG [1] compression standard is widely adopted in still images coding. The process consists of three stages: the block DCT transform, uniform quantization with a quantization table, and the Huffman (used in the baseline system) or the arithmetic (used in the extended system) entropy coder. In JPEG, one DCT coefficient has to be completely encoded before the coding of the next coefficient. More recently, a concept known as embedded or layered coding was proposed by Shapiro [3] and further developed by Taubman and Zakhor [4] in the context of wavelet coding. In contrast with the coefficient-by-coefficient approach, they adopted a new approach, in which each coefficient is successively quantized into a certain number of bits. The most significant bits of all coefficients are grouped together to form one layer and encoded first. Then, we move to the layer of the second significant bits and so on. The coding order is consistent with the importance of each bit so that the encoder and the decoder can stop at any time. The embedding property is essential for progressive image transmission. It also greatly simplifies the rate-control problem and allows unequal error protection for robust image transmission.

In this research, we generalize the layered coding to image compression methods based on the block DCT transform. Even though the generalization is not difficult, we feel that it is valuable to have its detailed implementation available. This is the main objective of the work. Besides providing more functionalities, the new coder gives a substantially better rate-distortion performance than the JPEG standard especially at low bit rates.

The work is organized as follows. We introduce the concept of layered coding in Section II. The detailed implementation is presented in Section III. Experimental results are given in Section IV.

2. OVERVIEW OF LAYERED CODING OF DCT COEFFICIENTS
The proposed layered DCT coding adopts the 8 x 8 block DCT transform which is the same as JPEG. The main differences lie in the quantization and the entropy coding schemes. For a given 8 x 8 DCT block, we arrange the 63 AC coefficients in a zig-zag order and denote them by $C_1, C_2, \ldots, C_{63}$. Suppose that the AC coefficients after quantization take a value ranging from $-32767$ to $32767$ so that each of them requires 16 bits for representation (including the sign bit). They are labeled with $S, B_1, \ldots, B_{15}$, where $S$ is the sign bit, $B_1$ the most significant bit (MSB) and $B_{15}$ the least significant bit (LSB). Values in such a bit matrix form a set of intermediate symbols. JPEG encodes these symbols in a coefficient-by-coefficient manner, with a 2-D run-level Huffman coder adopted in the JPEG baseline system while an arithmetic coder adopted in the extended system. The JPEG encoding scheme is illustrated with Fig. 1. First, all bits of coefficient $C_1$ are encoded, then all bits for coefficient $C_2$, and then $C_3$ and so on. The arithmetic coder in the JPEG extended
system adopts a two-way scan. The first scan starting from LSB $B_{15}$ to MSB $B_{6}$ locates the most significant bit $B_i$ for the current coefficient $C_j$. The second scan records bits $B_{i}, B_{i+1}, \ldots, B_{15}$ and the sign bit $S$. For more details of JPEG, we refer to [1]. One major issue with JPEG is the rate control problem. It is in general difficult to estimate the number of coding bits generated for a given Q factor. Note also that truncating the compressed bit stream at an arbitrary point is equivalent to deleting the bottom portion of the image, since it is encoded sequentially from one block to another.

In this research, we demonstrate that by adopting a different coding order for intermediate symbols, we convert the conventional JPEG coder into an embedded coder called the layered DCT coder. Consider the grouping of all bits $B_i$ of coefficients $C_j$, $1 \leq j \leq 63$, into layer $L_i$. Now, the new coder first encodes the most significant layer $L_1$, then layers $L_2, L_3$ and so on. Within each layer $L_i$, the coding follows the coefficient order, i.e. starting with coefficient $C_1$, then $C_2, C_3, \ldots, C_{63}$. This coding order is illustrated in Fig. 2. One important advantage of the layer-by-layer coding is that the resulting bit stream has the embedding property. Since the output bit stream is organized in an order of decreasing importance, rate-control can be easily achieved by simply truncating the bit stream according to the desired coding budget, or the desired coding quality. The layered DCT coding is also suitable for progressive transmission for the more important coding bit is always transmitted prior to the less important bit. Another advantage of the layer-by-layer coding is that the new scheme provides a better rate-distortion trade-off especially at low bit rates, which will be demonstrated in section IV.

3. DETAILED IMPLEMENTATION

We describe the detailed implementation of the proposed layered DCT compression algorithm in this section. The overall framework of our algorithm is depicted in Fig. 3. The distinctive feature of the layered DCT coder is that the quantization and the entropy coding procedures are carried out successively.

Step 1: Block DCT transform and coefficient scaling

As done in JPEG, the input image is partitioned into $8 \times 8$ blocks, and the block DCT transform is applied to each block. The DCT coefficients $C'_j$ are scaled with a standard JPEG quantization table $Q$, i.e.

$$C_j = \frac{C'_j}{Q_j}, j = 0, \ldots, 63.$$  

The scaling is performed to emphasize the visual importance of low frequency components.

Step 2: Coding of DC coefficients

DC coefficients of neighboring blocks are highly correlated. It can be either encoded with a differential layer coding [4] or encoded in the same way as the JPEG arithmetic coder. It is observed that the performance of the two schemes are very similar. In the experiment reported in Section IV, the latter approach is adopted.

Step 3: Successive quantization of AC coefficients

The main differences between the proposed method and JPEG are in the quantization scheme and the entropy coder. After the DCT transform and scaling,
JPEG applies an one-step quantization which maps each DCT coefficient to a value in a finite index set. The value is then converted to an intermediate symbol and encoded by an entropy coder. In the proposed new scheme, we adopt a successive quantization procedure which is achieved not in one step but with several successive steps. Roughly speaking, at layer \( i \), the DCT coefficient is only quantized up to the precision of the significant threshold \( T_i \). Then, the quantization result of layer \( i \) is refined with a smaller significant threshold \( T_{i+1} \) at layer \( i+1 \).

To initialize the process of successive quantization, we set all coefficients as insignificance and search the whole image for the maximum absolute value of the AC coefficients, which is denoted by \( T_0 \). Then, we apply the significant identification rule with significant threshold \( T_1 = T_0/2 \) to construct layer \( L_1 \). That is, for each scaled AC coefficient \( C_j \), if its magnitude is greater than \( T_1 \), we use symbol 1 to encode its significance, and then record its sign \( S_j \). Otherwise, we generate symbol 0 to indicate that it is still insignificant. Note that only the significant symbol needs a sign. In terms of mathematics, we have

\[
    C_j > T_1, \quad B_{j,1} = 1, \quad S_j = +, \quad E_{j,1} = C_j - 1.5T_1, \\
    C_j < -T_1, \quad B_{j,1} = 1, \quad S_j = -, \quad E_{j,1} = -C_j - 1.5T_1, \\
    \text{otherwise}, \quad B_{j,1} = 0, \quad E_{j,1} = C_j.
\]

where symbol \( E_{j,1} \) in the last column denotes the quantization residue at layer \( L_1 \). For each advanced layer \( L_{i+1}, i = 1, 2, \ldots \), we refine the significant threshold by half, i.e. \( T_{i+1} = T_i/2 \), and quantize all AC coefficients accordingly. If the AC coefficient to be coded is insignificant in all previous layers, the significance identification rule is applied. Otherwise, the refinement quantization rule is applied. These two rules can be summarized as follows:

1. rule of significance identification:

\[
    E_{j,i-1} > T_i, \quad B_{j,i} = 1, \quad S_j = +, \quad E_{j,i} = E_{j,i-1} - 1.5T_i, \\
    E_{j,i-1} < -T_i, \quad B_{j,i} = 1, \quad S_j = -, \quad E_{j,i} = -E_{j,i-1} - 1.5T_i, \\
    \text{otherwise}, \quad B_{j,i} = 0, \quad E_{j,i} = E_{j,i-1}.
\]

2. rule of refinement quantization:

\[
    E_{j,i-1} \geq 0, \quad B_{j,i} = 1, \quad E_{j,i} = E_{j,i-1} - T_i, \\
    E_{j,i-1} < 0, \quad B_{j,i} = 0, \quad E_{j,i} = -E_{j,i-1} - T_i.
\]

Figure 4: Illustration of significant symbol coding context, where * is the current coding position, o is the significant symbol of current layer and • is the significant symbol in the previous layer.

The generated symbols \( B_j \) and \( S_j \) provide a binary representation of the coefficient \( C_j \) normalized by \( T_0 \) as shown in Fig. 2. These symbols are generated with a decreasing order of importance and will be coded layer by layer in the next step.

Step 4: Context adaptive arithmetic coding

All symbols generated in successive quantization is binary, which can be coded efficiently by an arithmetic coder. Besides, we can predict the location of the significant coefficient with a context adaptive arithmetic coder used in the JPEG extended system. We also encode the refinement symbol and the sign with a specific context. The context adaptive arithmetic coder is a highly efficient entropy coder. Its average coding rate is close to the entropy of the source with low computational complexity. Its implementation only requires addition and shifting operations. There is no need for training or assuming the initial probability distribution so that it is parameter free. The source probability distributions \( p_0 \) and \( p_1 \) are estimated on the fly and implemented with a look-up table. The source probability distribution is represented as an 8 bit (1 byte) status of the coder, which includes 7 bits for the probability table and 1 bit for the most frequently appeared symbol (MFS). The small coder status enables the construction of a parallel coder for a compound source. We refer to [2] for more details.

For significance identification, our coding context
Table 1: Performance Comparison for Lena and Boat.

<table>
<thead>
<tr>
<th>Image</th>
<th>PSNR</th>
<th>Rate of JPEG</th>
<th>Layered DCT</th>
</tr>
</thead>
<tbody>
<tr>
<td>Lena</td>
<td>24.23</td>
<td>0.125</td>
<td>0.064</td>
</tr>
<tr>
<td></td>
<td>30.34</td>
<td>0.25</td>
<td>0.198</td>
</tr>
<tr>
<td></td>
<td>34.76</td>
<td>0.5</td>
<td>0.463</td>
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<tr>
<td></td>
<td>37.92</td>
<td>1.0</td>
<td>0.954</td>
</tr>
<tr>
<td>Boat</td>
<td>24.26</td>
<td>0.15</td>
<td>0.078</td>
</tr>
<tr>
<td></td>
<td>27.54</td>
<td>0.25</td>
<td>0.184</td>
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<tr>
<td></td>
<td>31.02</td>
<td>0.5</td>
<td>0.437</td>
</tr>
<tr>
<td></td>
<td>34.56</td>
<td>1.0</td>
<td>0.931</td>
</tr>
</tbody>
</table>

Figure 5: Rate-distortion performance of Lena using (a) JPEG Huffman coder (dashed line), (b) JPEG arithmetic coder (dotted line) and (c) the proposed layered DCT coder (solid line).

tradeoff curves for Lena are also depicted in Fig. 5.

We see from the experimental results that the layered DCT coder significantly outperforms the JPEG Huffman coder and also outperforms the JPEG arithmetic coder in most cases. The improvement is more substantial when the bit rates become lower. For example, it outperforms the JPEG Huffman and arithmetic coders by about 60% and 20%, respectively, when the bit rate is around 0.1 bpp. In addition to the superior rate-distortion performance, the layered DCT coder possesses the embedding property which makes progressive image transmission and rate-control easier to attain.

4. EXPERIMENTAL RESULTS

Experiments are conducted to compare the new layered DCT coder with the JPEG standard using the Huffman and the arithmetic entropy coders. The images used in the experiments are Lena and Boat of size 512 x 512. For fair comparison, we strip the header of the JPEG coded bit stream. We use the same quantization factor Q for JPEG Huffman and arithmetic coders, which results in the same coding image quality for both coders. Due to the embedding property of the proposed layered DCT coder, we can truncate the coded bit stream when it reaches the same PSNR values as JPEG. By doing so, we can compare the bit rates for the three coding schemes with the same quality. The results are shown in Table 1. We list the coding bit rate reduction with respect to JPEG Huffman and arithmetic coders in the last two columns of the table. The rate-distortion tradeoff curves for Lena are also depicted in Fig. 5.

5. REFERENCES